

Weak Hiding for C++ Concepts and a Generic Way to Understand Name Binding

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Objectives

- ❖ A simpler way to describe name binding.
- ❖ A more powerful way to describe name binding.
- ❖ Applying this description to C++ concepts.



Our Tools

- ❖ scope combinators,
- ❖ an abstraction of programming languages, and
- ❖ a distinction between **name lookup** and **name resolution**.



Outline

① Current name binding mechanisms:

1. Across languages.
2. Specific to C++:
 - a. Argument-dependent lookup (ADL).
 - b. Uses of operators.

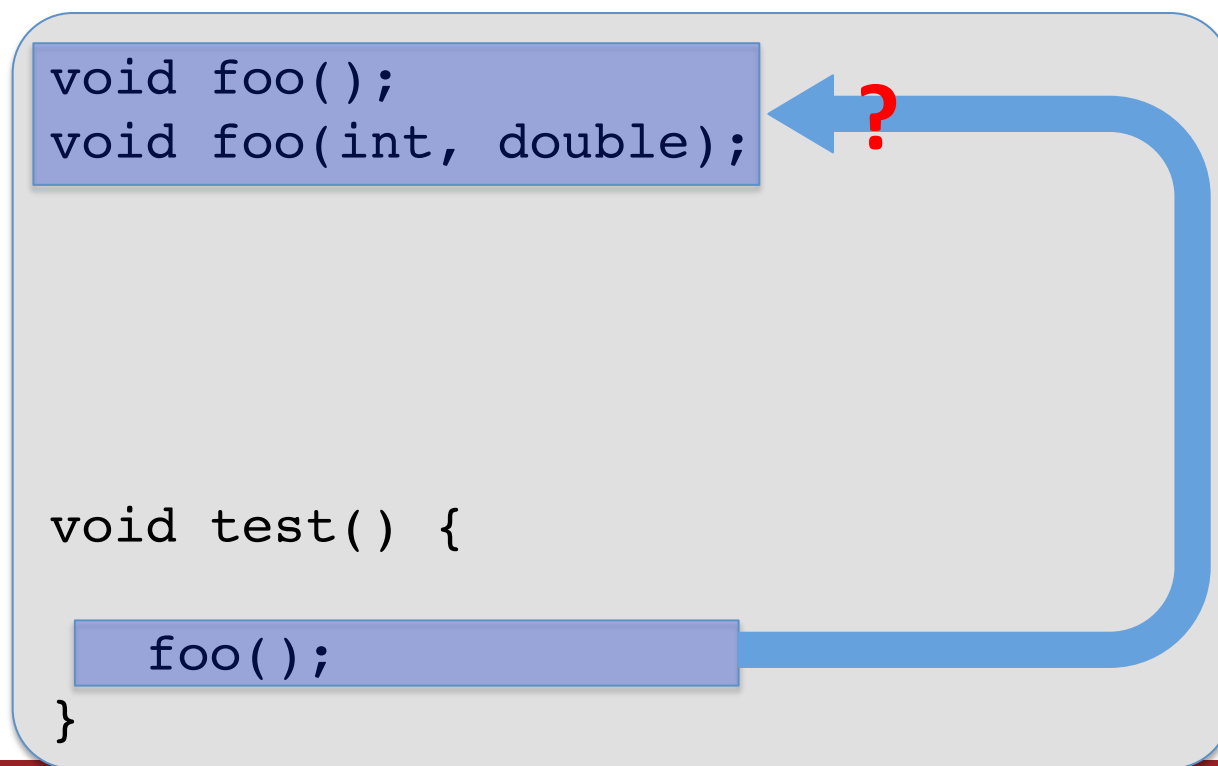
② A new scoping rule: **Weak hiding**

- Implementation: **two-stage name binding (Bind^{x2})**
- Application: Preventing code breakage when switching to concepts-enabled libraries.



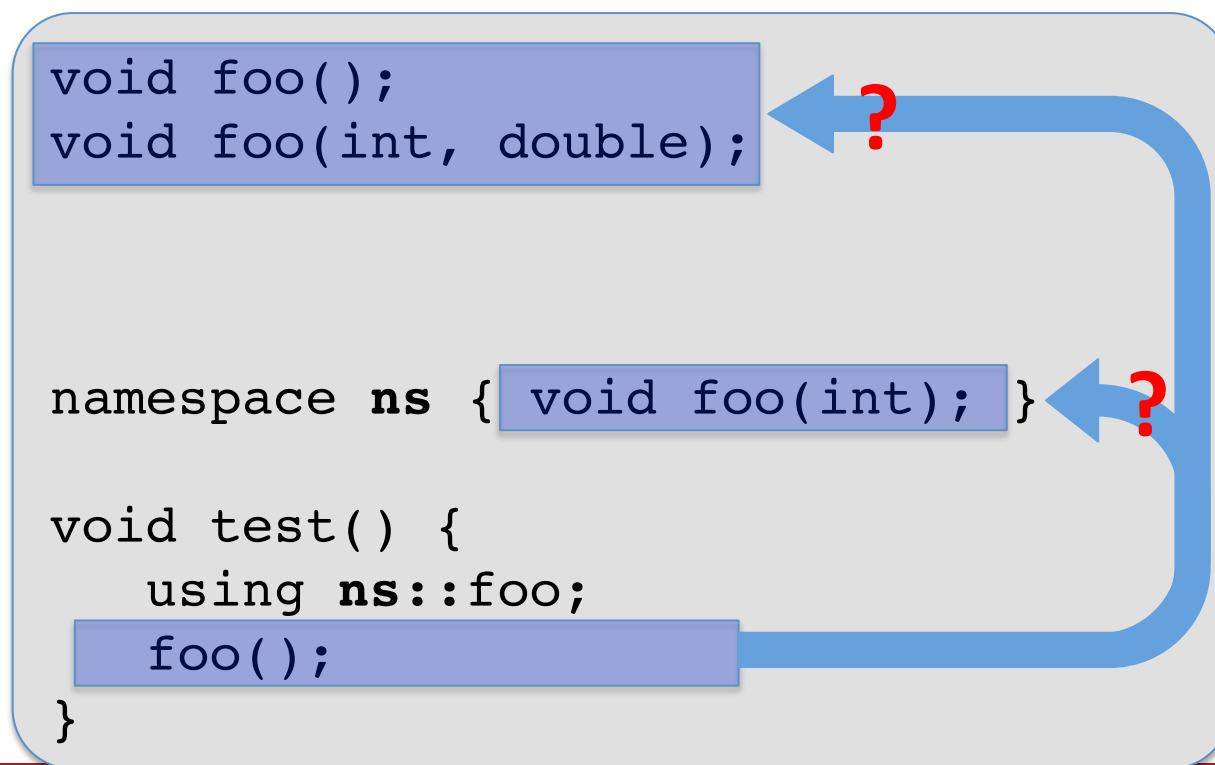
Name Binding

- ❖ Deciding what a particular use of a name refers to.
- ❖ Use of a name = **a *reference***; What it refers to = **a *declaration***.



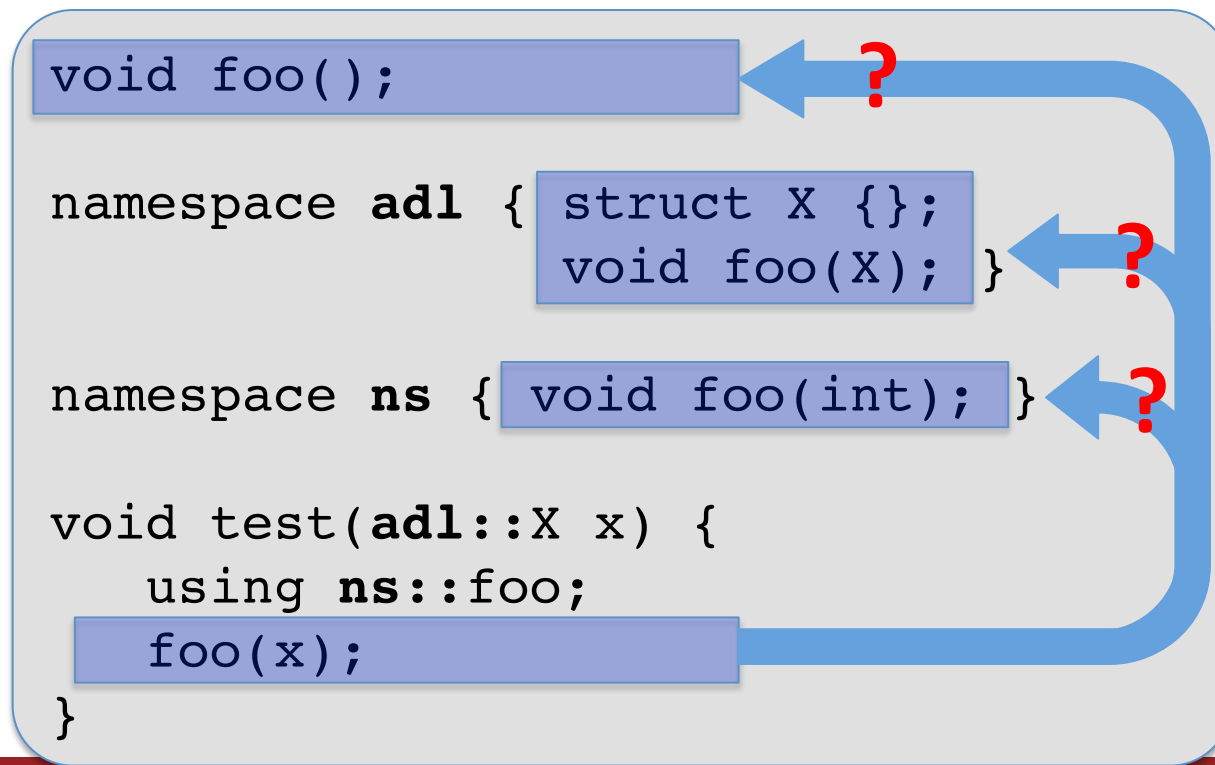
Name Binding

❖ Binds a *reference* to a *declaration*.



Name Binding

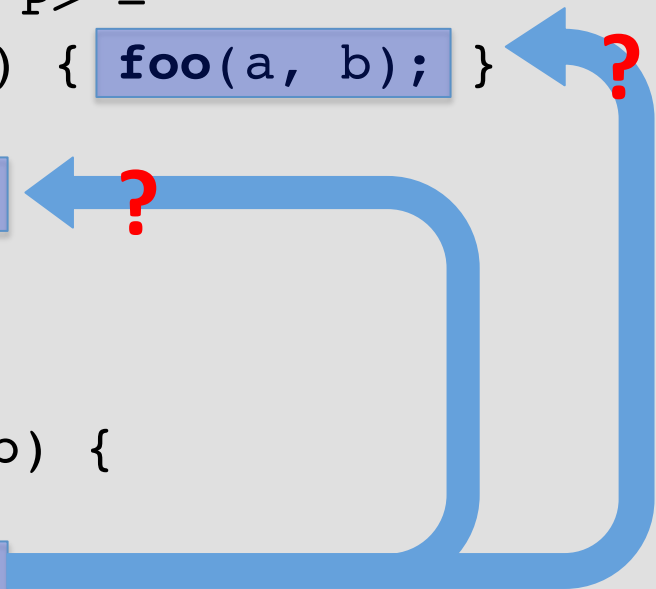
❖ Binds a *reference* to a *declaration*.



Name Binding

❖ Binds a *reference* to a *declaration*.

```
concept Foo2<typename P> =  
    requires (P a, P b) { foo(a, b); }
```



```
void foo(int) { }
```

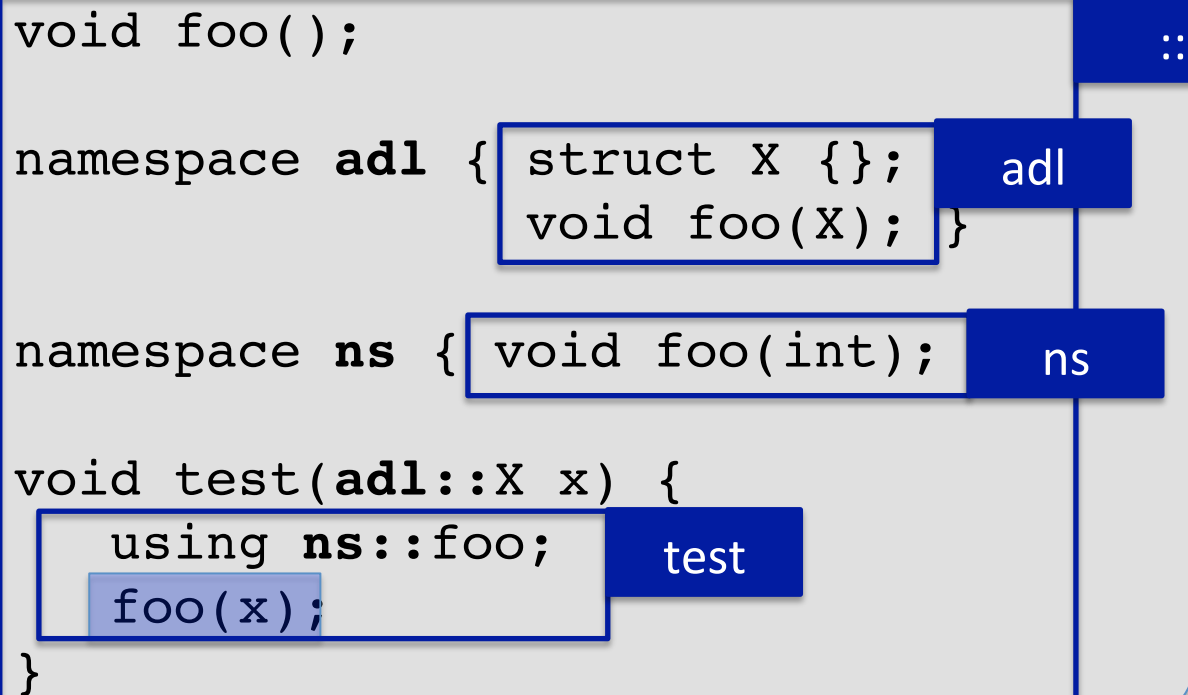
```
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```



Name Binding

- ❖ Depends on a **scope**, nested or combined with other scopes.

```
void foo();  
  
namespace adl { struct X {}; void foo(X); }  
  
namespace ns { void foo(int); }  
  
void test(adl::X x) {  
    using ns::foo;  
    foo(x);  
}
```



Name Binding

❖ Differs between **languages, designs, or kinds of references:**

- argument-dependent lookup (ADL),
- simple function calls,
- uses of operators,
- uses of types,
- C++ Multimethods,
- C++,
- Haskell,
- etc...

```
void foo();

namespace adl { struct X {}; void foo(X); }

namespace ns { void foo(int); }

void test(adl::X x, adl::X y) {
    using ns::foo;
    foo(x + y);
}
```



Standards for Name Binding

- ❖ Can be challenging to understand, analyze, and compare.
- ❖ Are intertwined with other sections of the standard.

▪ C++ Name Binding Specification, at a glance:

- | | |
|--------------------------|---------------------------|
| ▪ 3.4 – Name lookup | 13 pages, 45+ paragraphs |
| ▪ 14.6 – Name resolution | 13 pages, 50+ paragraphs |
| ▪ 13 – Overloading | 30 pages, 140+ paragraphs |
| ▪ 11 – Member access | 10 pages, 35+ paragraphs |



Standards for Name Binding

- ❖ Typically depend on details of languages and the varied features they support.

ConceptC++:

```
template<Foo2 T>
void test(T) {
    foo();
}
```

- ⇒ C++ templates + Concepts
- ⇒ (Non)dependent names
- ⇒ Overload resolution
- ⇒ Namespaces + using declarations
- ⇒ C++ compilers, e.g., GCC, Clang, ...
- ⇒ Etc...

Haskell:

```
test::Foo2 t => t -> ()
test _ = foo()
```

- ⇒ Hindley-Milner type system + Type classes
- ⇒ Type inference + Dependency analysis
- ⇒ Dictionary-passing style
- ⇒ Modules and import declarations
- ⇒ Haskell compilers, e.g., GHC, Hugs, NHC, ...
- ⇒ Etc...



Standards for Name Binding

❖ Are understood differently by different compilers.

```
namespace adl { struct X {};  
                void foo(X); }  
  
namespace ns { void foo(int); }  
  
void bar(adl::X x) {  
    using ns::foo;  
    foo(x);  
}  
  
void baz(adl::X x) {  
    void foo(int);  
    foo(x);  
}
```

In bar():

Success. ADL enabled.
Binds **foo(x)** to **adl::foo()**.

In baz():

GCC, Clang, Comeau:

Failure. ADL Disabled.
foo(x) does not match **baz::foo(int)**.

Intel:

Success.



Standards for Name Binding

❖ Are understood differently by different compilers.

```
struct X {}; struct Y {};
```

```
void operator+(X, X) { }  
void operator+(X, Y) { }
```

```
void test(X x, Y y) {  
    void operator+(X, X);  
    x + x;  
    x + y;  
    operator+(x, x);  
    operator+(x, y);  
}
```

GCC, Clang, MS Visual Studio editor :

x + y succeeds.

operator+(x, y) fails.

Intel, MS Visual Studio:

x + y succeeds.

operator+(x, y) succeeds.

Comeau:

x + y fails.

operator+(x, y) fails.



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Standards for Name Binding

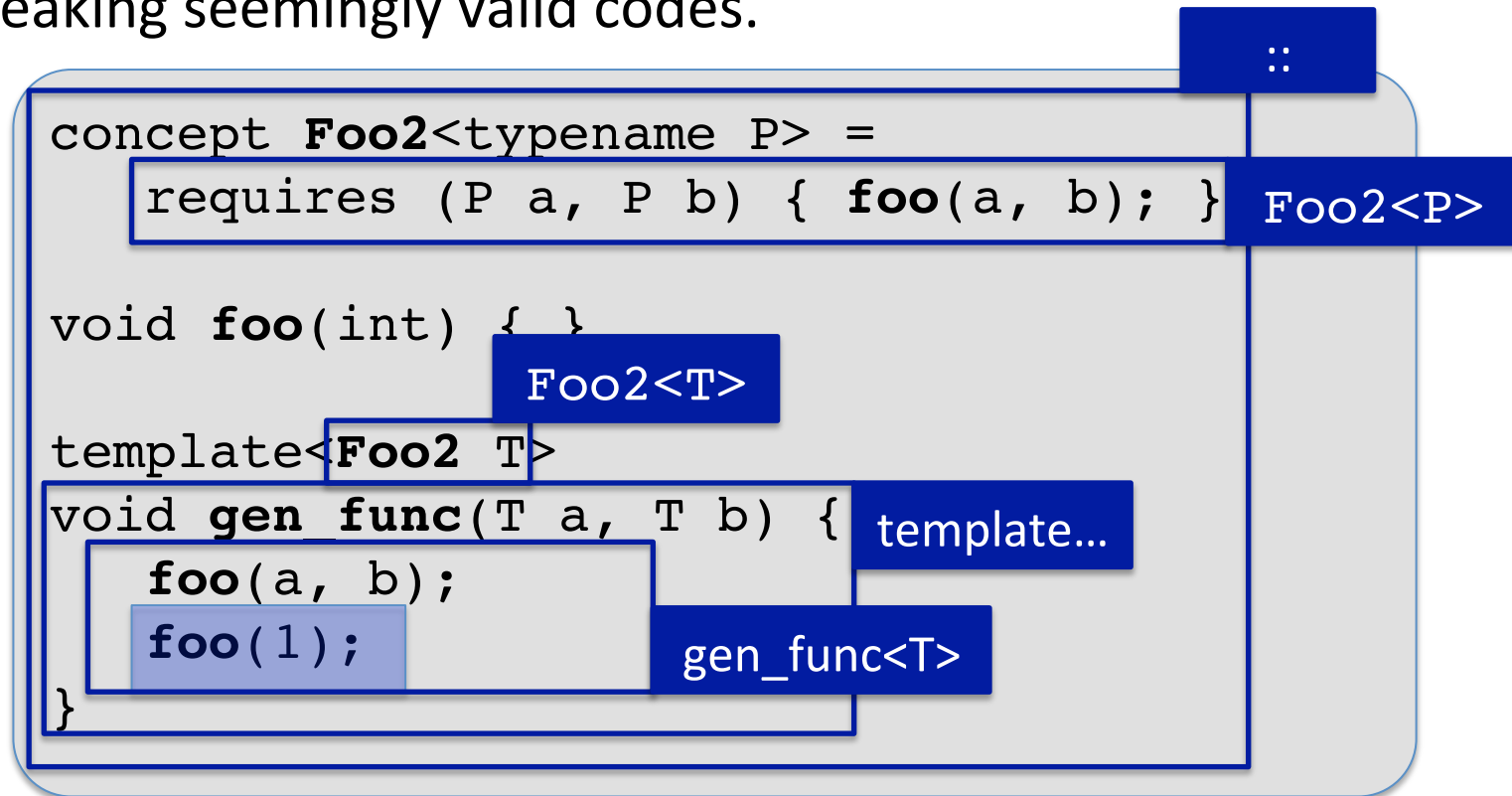
- ❖ Gain complexity with the addition of concepts,
- ❖ breaking seemingly valid codes.

```
concept Foo2<typename T> =  
    requires (T a, T b) { foo(a, b); }  
  
void foo(int) { }  
  
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```



Name Binding Specifications

- ❖ Gain complexity with the addition of concepts,
- ❖ breaking seemingly valid codes.



Name Binding Specifications

- ❖ Gain complexity with the addition of concepts,
- ❖ breaking seemingly valid codes.

The diagram shows a C++ code snippet with several annotations and callouts:

```
concept Foo2<typename P> =  
    requires (P a, P b) { foo(a, b); }  
  
void foo(int) { }  
  
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```

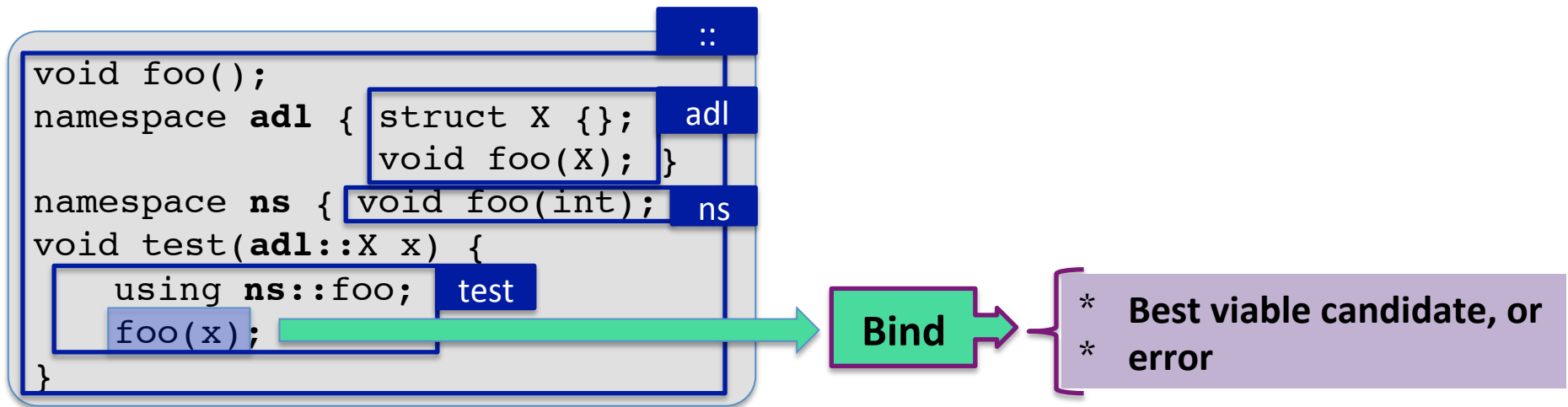
Annotations and Callouts:

- A blue box labeled `Foo2<T>` points to the `Foo2` template parameter in `template<Foo2 T>`.
- A blue box labeled `::` points to the namespace separator in the concept definition.
- A red callout bubble points to the `foo` call inside the concept's `requires` clause: "Should `Foo2<T>::foo()` shadow `::foo()`? How?"
- A red callout bubble points to the `foo(1);` call inside `gen_func`: "Reject or accept the call `foo(1)`?"



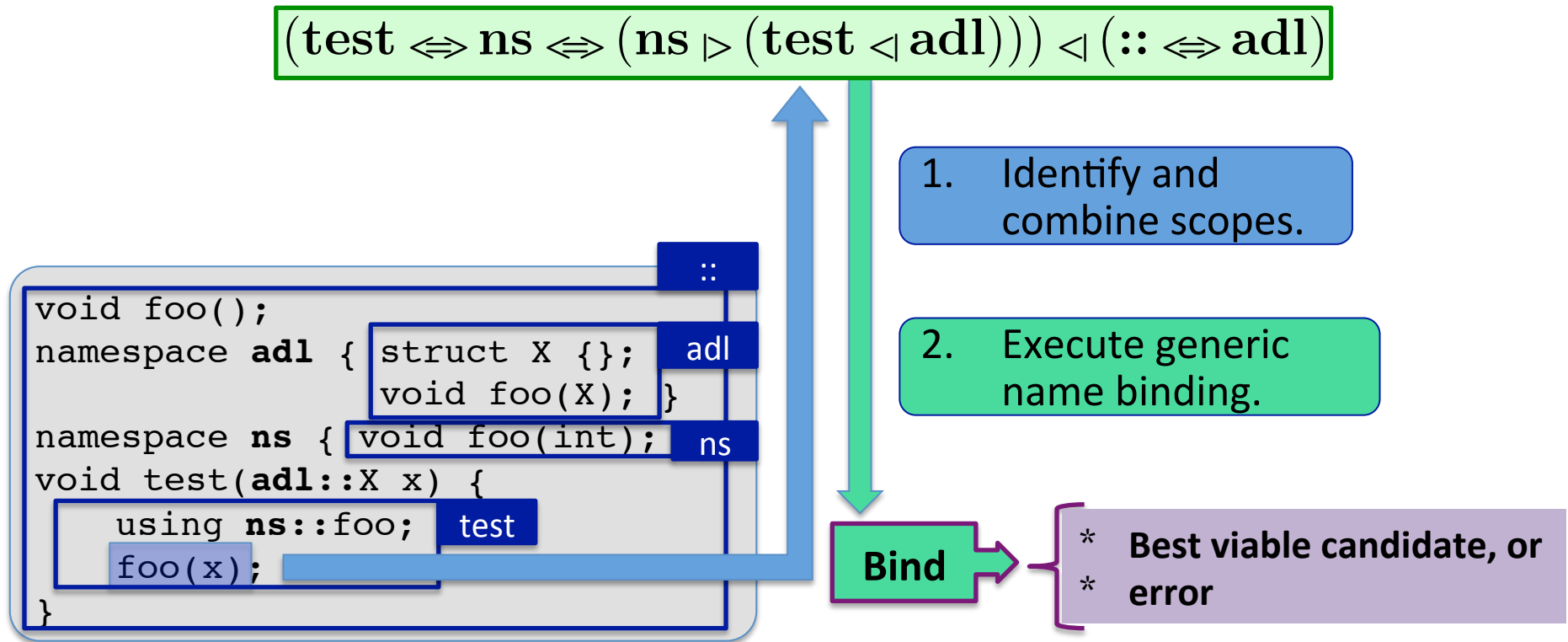
Our Name Binding Framework

❖ Specifies name binding.



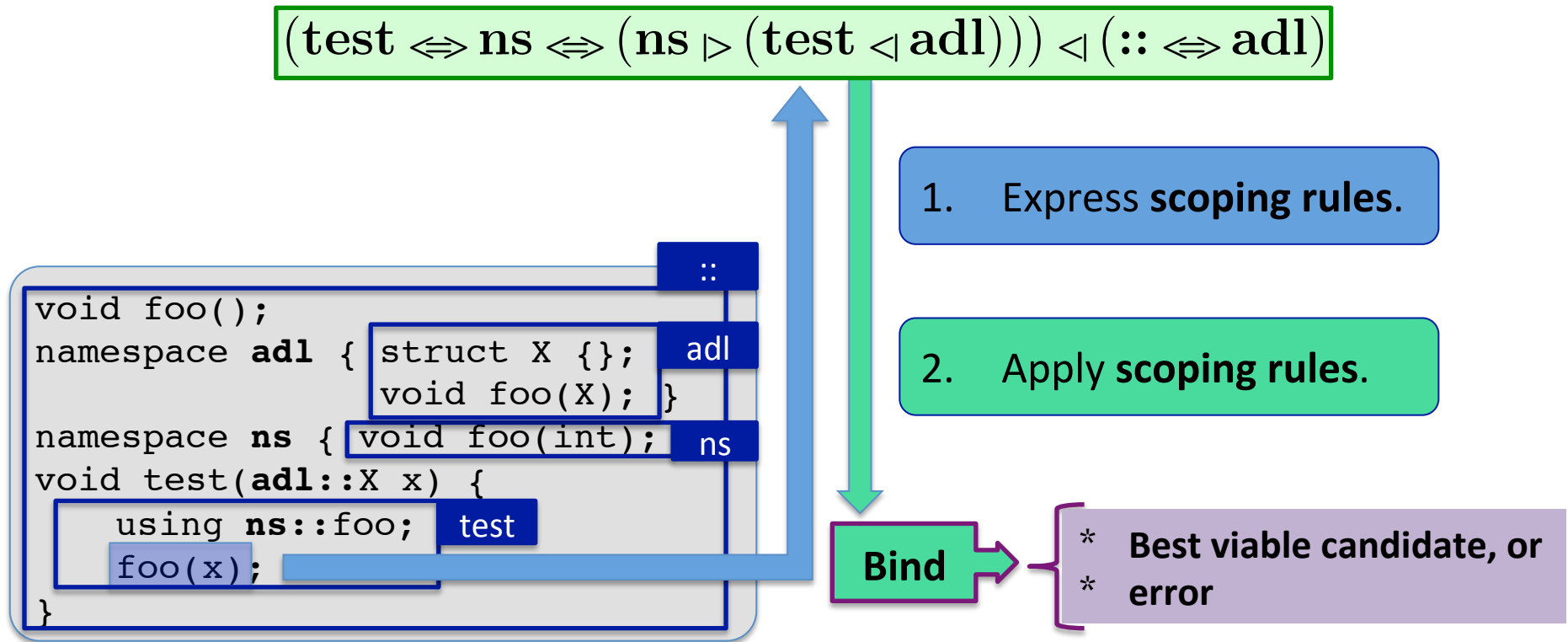
Our Name Binding Framework

- ❖ Specifies name binding, independently of the language.



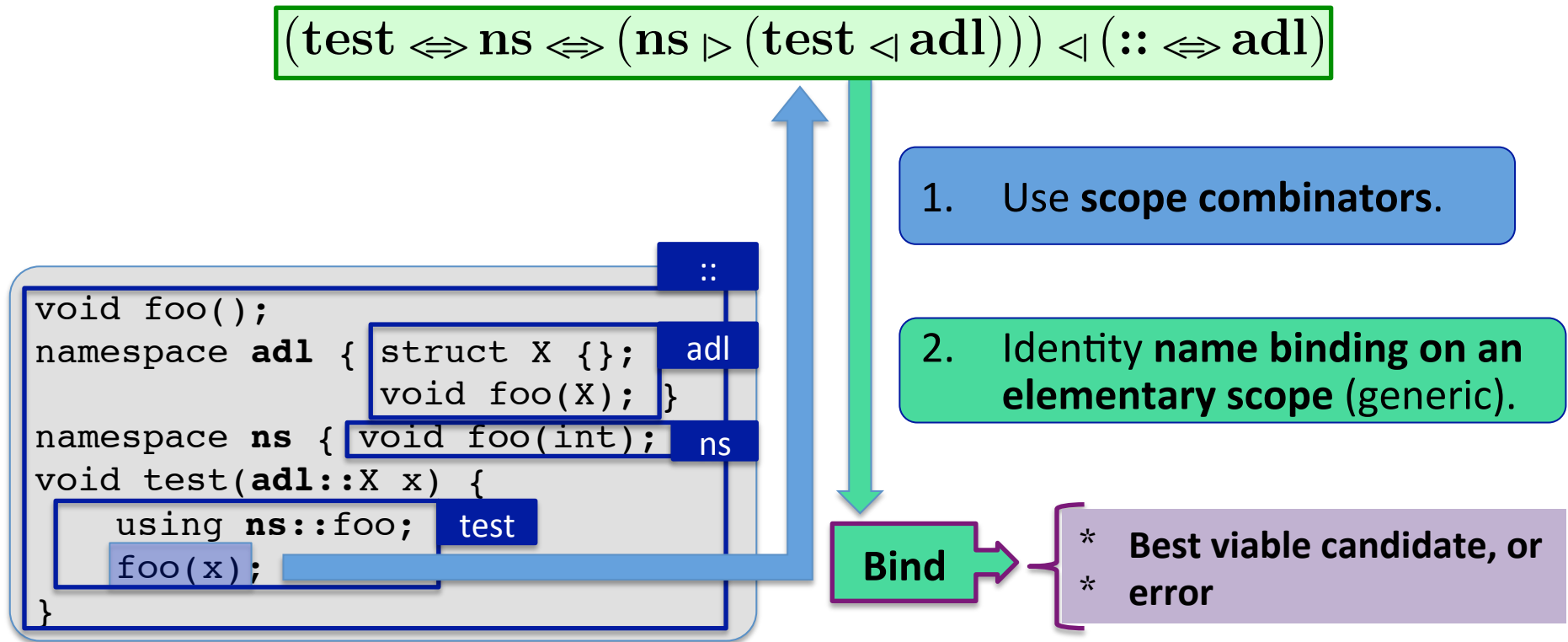
Our Name Binding Framework

- ❖ Specifies name binding, independently of the language.



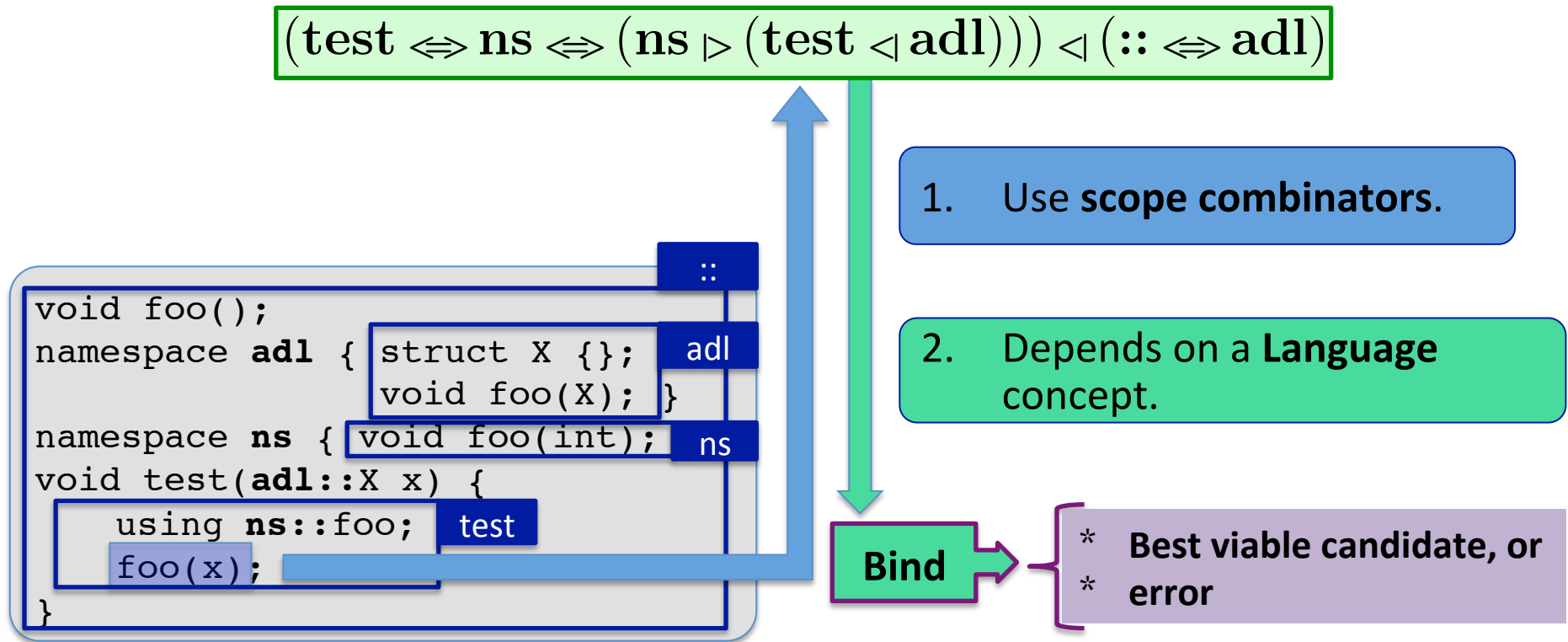
Our Name Binding Framework

- ❖ Specifies name binding, independently of the language.



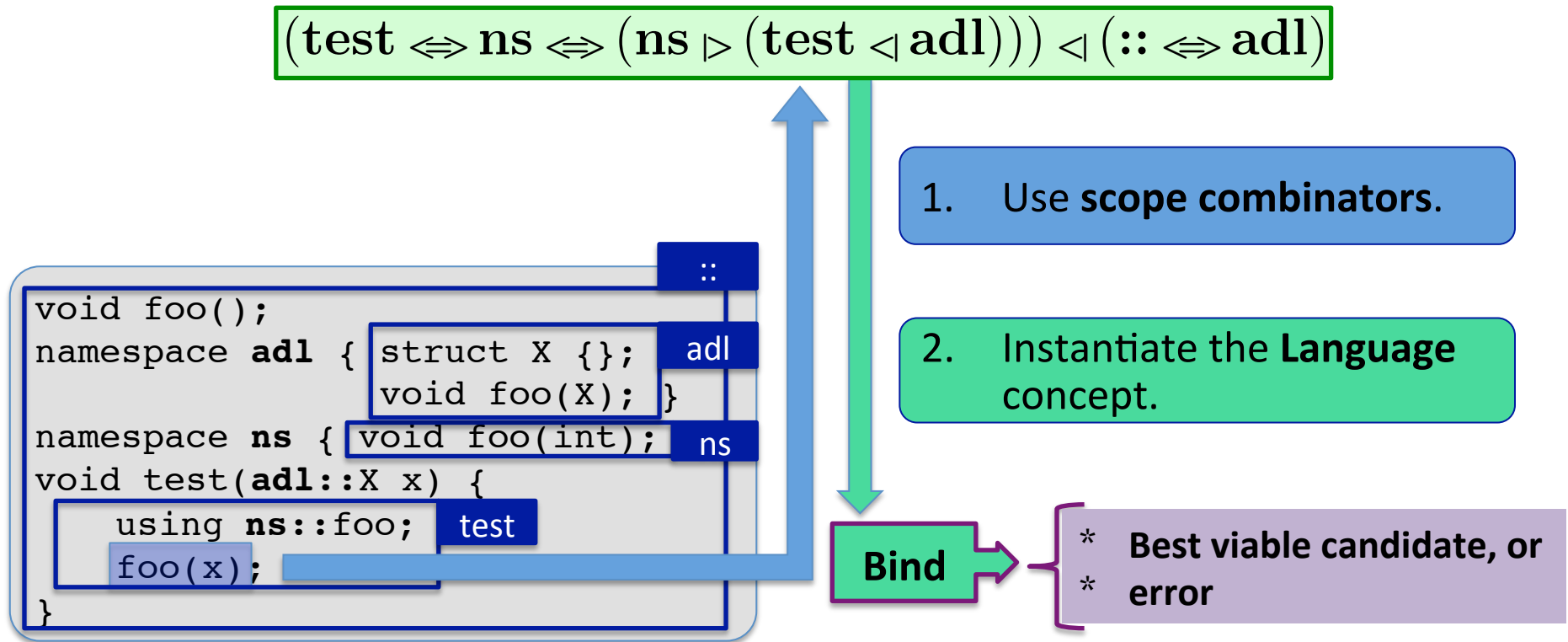
Our Name Binding Framework

- ❖ Specifies name binding, independently of the language.



Our Name Binding Framework

- ❖ Specifies name binding, independently of the language.



Our Name Binding Framework

❖ Abstracts from *declarations*, *references*, and *scopes*.

$$\text{bind} : \text{Ref} \times \text{Scope}_{\text{Ref}, \text{Decl}} \rightarrow (\text{Decl} + \text{Error})$$

```
void foo();  
namespace adl { struct X {}; void foo(X); }  
namespace ns { void foo(int); }  
void test(adl::X x) {  
    using ns::foo;  
    foo(x);  
}
```

Scopes as maps of **references** to **sets of matching declarations**.

Bind

* Best viable candidate, or
* error



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Our Name Binding Framework

- ❖ Views name binding as composed of **name lookup** and **resolution**.

$$\text{bind} : \text{Ref} \times \text{Scope}_{\text{Ref}, \text{Decl}} \rightarrow (\text{Decl} + \text{Error})$$
$$\text{bind}(\text{ref}, \text{scope}) = ((\text{resolve } \text{ref}) \circ (\text{lookup } \text{ref})) \text{ scope}$$

```
void foo();  
namespace adl { struct X {}; void foo(X); }  
namespace ns { void foo(int); }  
void test(adl::X x) {  
    using ns::foo;  
    foo(x);  
}
```

Name lookup returns the **set of matching declarations**, for a given reference.

Bind

* Best viable candidate, or
* error



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1. Expressing Scoping Rules

$$(\text{test} \Leftrightarrow \text{ns} \Leftrightarrow (\text{ns} \triangleright (\text{test} \triangleleft \text{adl}))) \triangleleft (:: \Leftrightarrow \text{adl})$$

1. Use **scope combinators**.

2. Instantiate the **Language** concept.

Bind

* Best viable candidate, or
* error

$$\text{Scope}_{\text{Ref,Decl}} \times \text{Scope}_{\text{Ref,Decl}} \rightarrow \text{Scope}_{\text{Ref,Decl}}$$

```

void foo();
namespace adl { struct X {}; void foo(X); }
namespace ns { void foo(int); }
void test(adl::X x) {
    using ns::foo;
    foo(x);
}
    
```



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The Combinators

❖ *Hiding*: \triangleleft

- Commonly known as “shadowing”.

❖ *Merging*: \Leftrightarrow

- Usually the alternative option to “shadowing”.

❖ *Opening*: \triangleright

- [New name] Necessary to describe ADL.
- A dual of *hiding*.

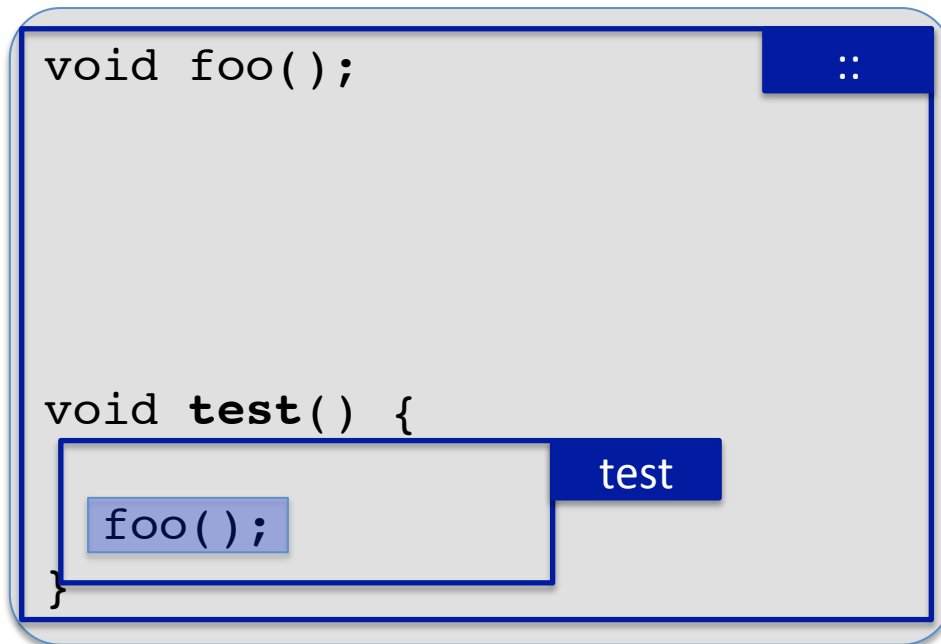
❖ *Weak Hiding*: \Leftarrow

- [New rule] Necessary for (C++) concepts.
- A sweet middle between *hiding* and merging.



The *Hiding* Combinator (\triangleleft)

$$s_1 \triangleleft s_2 = \text{lookup}_{ref} s_1 ? \text{lookup}_{ref} s_1 : \text{lookup}_{ref} s_2$$



`test \triangleleft ::`

Result:

Binds **foo()** to **::foo()**.



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The *Merging* Combinator ($\Leftarrow \Rightarrow$)

$$s_1 \Leftarrow \Rightarrow s_2 = lookup_{ref} s_1 \cup lookup_{ref} s_2$$

```
void foo();
```

```
::
```

```
namespace ns {
```

```
void foo(int);
```

```
ns
```

```
}
```

```
void test() {
```

```
using ns::foo;
```

```
test
```

```
foo();
```

```
}
```

$(test \Leftarrow \Rightarrow ns) \triangleleft ::$

Result:

Finds **ns::foo()**;

Fails to bind **foo()**.



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The *Weak Hiding* Combinator (\Leftarrow)

$$s_1 \Leftarrow s_2 = (\text{resolve}_{ref} \circ \text{lookup}_{ref}) s_1 ?$$

$$\text{lookup}_{ref} s_1 : \text{lookup}_{ref} s_2$$

```
void foo();
```

```
::
```

```
namespace ns {
```

```
void foo(int);
```

```
ns
```

```
}
```

```
void test() {
```

```
using ns::foo;
```

```
test
```

```
foo();
```

```
}
```

```
(test  $\Leftarrow$  ns)  $\Leftarrow$  ::
```

Result:

Binds **foo()** to **::foo()**.



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The *Opening* Combinator (\triangleright)

$$s_1 \triangleright s_2 = \text{lookup}_{ref} s_1 ? \text{lookup}_{ref} s_2 : \text{empty}$$

```
void foo();  
namespace ns { void foo(int); }  
namespace adl { struct X {}; void foo(typ); }  
void test(adl::X x) {  
    using ns::foo;  
    foo(x);  
}
```

Diagram illustrating the scope resolution process for the `foo(x);` call in the `test` function. The `test` function is in the global namespace. It uses `ns::foo`. The `ns` namespace is defined in the global namespace. The `adl` namespace is defined in the global namespace. The `test` function is in the global namespace.

$$(\text{test} \Leftrightarrow \text{ns} \Leftrightarrow (\text{ns} \triangleright (\text{test} \triangleleft \text{adl}))) \\ \triangleleft (:\Leftrightarrow \text{adl})$$

Result:

Finds `ns::foo()`;

Enables ADL;

Binds `foo(x)` to `adl::foo()`.



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The *Opening* Combinator (\triangleright)

$$s_1 \triangleright s_2 = \text{lookup}_{ref} s_1 ? \text{lookup}_{ref} s_2 : \text{empty}$$

```
void foo();  
namespace ns { void foo(int); }  
namespace adl { struct X {}; void foo(typ); }  
void test(adl::X x) {  
    using ns::foo;  
    void foo();  
    foo(x);  
}
```

Diagram illustrating the scope resolution for the `test` function. The `test` function is defined in the `adl` namespace. It uses `using ns::foo;` to bring the `foo` function from the `ns` namespace into the `test` function's scope. The `foo(x);` call is then resolved to `ns::foo` because `ns::foo` is in the current scope.

$$(\text{test} \Leftrightarrow \text{ns} \Leftrightarrow (\text{ns} \triangleright (\text{test} \triangleleft \text{adl}))) \\ \triangleleft (:: \Leftrightarrow \text{adl})$$

Result:

Finds **`test::foo();`**

Disables ADL;

Fails to bind **`foo(x).`**



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The Combinators – Recap

$\langle |, \Leftrightarrow, | \rangle, \Leftarrow$ $\text{Scope}_{\text{Ref,Decl}} \times \text{Scope}_{\text{Ref,Decl}} \rightarrow \text{Scope}_{\text{Ref,Decl}}$

$$s_1 \Leftrightarrow s_2 = \text{lookup}_{\text{ref}} s_1 \cup \text{lookup}_{\text{ref}} s_2$$

$$s_1 \langle | s_2 = \text{lookup}_{\text{ref}} s_1 ? \text{lookup}_{\text{ref}} s_1 : \text{lookup}_{\text{ref}} s_2$$

$$s_1 | \rangle s_2 = \text{lookup}_{\text{ref}} s_1 ? \text{lookup}_{\text{ref}} s_2 : \text{empty}$$

$$s_1 \Leftarrow s_2 = (\text{resolve}_{\text{ref}} \circ \text{lookup}_{\text{ref}}) s_1 ?$$

$$\text{lookup}_{\text{ref}} s_1 : \text{lookup}_{\text{ref}} s_2$$

Special cases of a “conditional” combinator.



Applications

❖ Understanding current name binding mechanisms:

- Argument-dependent lookup (ADL).
- C++ operators.
- A cross-language analysis.

❖ Exploring concepts designs

- Understanding the current limitations.
- Exploring new solutions:
 - weak hiding,
 - 2-Stage name binding, and
 - parameterized weak hiding.

❖ Simplifying compiler designs.



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ADL Example

```
namespace ns1 { struct X {};  
                ns1 void foo(X); }  
namespace ns2 { void foo(int); }  
void bar(ns1::X x) {  
    using ns2::foo; bar  
    foo(x);  
}  
void baz(ns1::X x) {  
    void foo(int); baz  
    foo(x);  
    {  
        using ns2::foo; // H  
        foo(x); H  
    }  
}
```

In bar():

In baz():

In baz()'s inner scope, H:



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ADL Example

```
namespace ns1 { struct X {};  
                 ns1 void foo(X); }  
namespace ns2 { void foo(int); }  
void bar(ns1::X x) { ns2  
    using ns2::foo; bar  
    foo(x);  
}  
void baz(ns1::X x) {  
    void foo(int);  
    foo(x);  
    { // H  
        using ns2::foo;  
        foo(x);  
    }  
}
```

In bar():

$$(bar \Leftrightarrow ns2 \Leftrightarrow (ns2 \triangleright (bar \triangleleft ns1))) \\ \triangleleft (:: \Leftrightarrow ns1)$$

In baz():

In baz()'s inner scope, H:



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ADL Example

```
namespace ns1 { struct X {};  
                 ns1 void foo(X); }  
namespace ns2 { void foo(int); }  
void bar(ns1::X x) {  
    using ns2::foo;  
    foo(x);  
}  
void baz(ns1::X x) {  
    void foo(int);  
    foo(x);  
    {  
        using ns2::foo;  
        foo(x);  
    }  
}
```

In bar():

$$(bar \Leftrightarrow ns2 \Leftrightarrow (ns2 \triangleright (bar \triangleleft ns1))) \\ \triangleleft (:: \Leftrightarrow ns1)$$

In baz():

$$baz \triangleleft (:: \Leftrightarrow ns1)$$

In baz()'s inner scope, H:



ADL Example

```

namespace ns1 { struct X {};
                 ns1 void foo(X); }
namespace ns2 { void foo(int); }
void bar(ns1::X x) {
    using ns2::foo;
    foo(x);
}
void baz(ns1::X x) {
    void foo(int);
    foo(x);
    {
        using ns2::foo;
        foo(x);
    }
}

```

Diagram illustrating the scope resolution and ADL (Argument-Dependent Lookup) process. The code defines two namespaces, `ns1` and `ns2`, and two functions, `bar` and `baz`. `bar` is defined in `ns1` and takes an argument of type `ns1::X`. `baz` is also defined in `ns1` and takes an argument of type `ns1::X`. Inside `baz`, there is an inner scope `H` where `using ns2::foo;` is used. The diagram highlights the scope resolution process, showing that `baz` is in `ns1`, and the inner scope `H` is also in `ns1`. The `using` statement in `H` brings `ns2::foo` into scope, and the `foo(x);` call in `H` is resolved using ADL.

In bar():

$$(\text{bar} \Leftarrow \text{ns2} \Leftarrow (\text{ns2} \triangleright (\text{bar} \triangleleft \text{ns1})))$$

$$\triangleleft (:: \Leftarrow \text{ns1})$$

In baz():

$$\text{baz} \triangleleft (:: \Leftarrow \text{ns1})$$

In baz()'s inner scope, H:

$$(\text{H} \Leftarrow \text{ns2} \Leftarrow (\text{ns2} \triangleright (\text{H} \triangleleft \text{ns1})))$$

$$\triangleleft \text{baz} \triangleleft (:: \Leftarrow \text{ns1})$$


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ADL Example

```

namespace ns1 { struct X {};
                 ns1 void foo(X); }
namespace ns2 { void foo(int); }
void bar(ns1::X x) {
    using ns2::foo;
    foo(x);
}
void baz(ns1::X x) {
    void foo(int);
    foo(x);
    {
        using ns2::foo;
        foo(x);
    }
    // H
}

```

In bar():

$f(\text{bar}) \triangleleft (:: \Leftrightarrow \text{ns1})$

In baz():

$f(\text{baz}) \triangleleft (:: \Leftrightarrow \text{ns1})$

In baz()'s inner scope, H:

$f(H) \triangleleft f(\text{baz}) \triangleleft (:: \Leftrightarrow \text{ns1})$

----- with -----

$$f(X) = X \Leftrightarrow U_X \Leftrightarrow (U_X \triangleright (X \triangleleft \text{ns1}))$$

$$U_{\text{bar}} = U_H = \text{ns2}$$

$$U_{\text{baz}} = \text{empty}$$



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ADL Scoping Rules

$$f_{ADL}(H) \triangleleft \langle \langle_{i=1}^s f_{ADL}(S_i) \triangleleft \langle \langle_{i=1}^{n-1} fn_{ADL}(N_i) \triangleleft (N_n \Leftrightarrow U_{N_n} \Leftrightarrow ((\tilde{N}_n \Leftrightarrow \tilde{U}_{N_n}) \triangleleft ADL)) \rangle \rangle$$

$$f_{ADL}(X) = X \Leftrightarrow U_X \Leftrightarrow (U_X \triangleright ((X \Leftrightarrow \tilde{U}_X) \triangleleft ADL)),$$

$S_1 \cdots S_s$ = surrounding non-namespace scopes,

$$fn_{ADL}(N) = N \Leftrightarrow U_N \Leftrightarrow ((N \Leftrightarrow U_N) \triangleright ((\tilde{N} \Leftrightarrow \tilde{U}_N) \triangleleft ADL)),$$

$N_1 \cdots N_n$ = surrounding namespaces,

H = scope where name binding is triggered from,

U_X = using declarations in scope X ,

\tilde{X} = non-function (template) declarations in scope X , and

ADL = associated namespaces of reference's arguments.



ADL Scoping Rules

when scope **H** is an inner namespace scope

$$\text{fn}_{\text{ADL}}(\mathbf{H}) \triangleleft \langle \rangle_{i=1}^{n-1} \text{fn}_{\text{ADL}}(\mathbf{N}_i) \triangleleft \left(\mathbf{N}_n \Leftrightarrow \mathbf{U}_{\mathbf{N}_n} \Leftrightarrow \left(\left(\tilde{\mathbf{N}}_n \Leftrightarrow \tilde{\mathbf{U}}_{\mathbf{N}_n} \right) \triangleleft \text{ADL} \right) \right)$$

$$\text{fn}_{\text{ADL}}(\mathbf{N}) = \mathbf{N} \Leftrightarrow \mathbf{U}_{\mathbf{N}} \Leftrightarrow \left(\left(\mathbf{N} \Leftrightarrow \mathbf{U}_{\mathbf{N}} \right) \triangleright \left(\left(\tilde{\mathbf{N}} \Leftrightarrow \tilde{\mathbf{U}}_{\mathbf{N}} \right) \triangleleft \text{ADL} \right) \right),$$

$\mathbf{N}_1 \cdots \mathbf{N}_n$ = surrounding namespaces,

\mathbf{H} = scope where name binding is triggered from,

$\mathbf{U}_{\mathbf{X}}$ = using declarations in scope \mathbf{X} ,

$\tilde{\mathbf{X}}$ = non-function (template) declarations in scope \mathbf{X} , and

ADL = associated namespaces of reference's arguments.



ADL Scoping Rules

when scope **H** is the outermost namespace scope

$$\left(N_n \Leftrightarrow U_{N_n} \Leftrightarrow \left(\left(\tilde{N}_n \Leftrightarrow \tilde{U}_{N_n} \right) \triangleleft \text{ADL} \right) \right)$$

$$H = N_n$$

H = scope where name binding is triggered from,

U_X = using declarations in scope **X**,

\tilde{X} = non-function (template) declarations in scope **X**, and

ADL = associated namespaces of reference's arguments.



C++ Operators Example

```
struct X {}; struct Y {};
```

```
void operator+(X, X) { }
```

```
void operator+(X, Y) { }
```

```
void test(X x, Y y) {
```

```
    void operator+(X, X); test
```

```
    x + x;
```

```
    x + y;
```

```
    operator+(x, x);
```

```
    operator+(x, y);
```

```
}
```

```
::
```

(test \Leftrightarrow builtin-ops \Leftrightarrow X) \triangleleft ::



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C++ Operators Scoping Rules

without ADL

$$\left(H \Leftrightarrow U_H \Leftrightarrow O_B \Leftrightarrow M \right) \\ \triangleleft \bigwedge_{i=1}^s S_i \triangleleft \bigwedge_{i=1}^n N_i$$

$S_1 \cdots S_s$ = surrounding non-namespace scopes,

$N_1 \cdots N_n$ = surrounding namespaces,

H = scope where name binding is triggered from,

O_B = built-in operators,

M = member scope of operator's first argument,

U_X = using declarations in scope X , and



C++ Operators Scoping Rules with ADL

$$\left(H \Leftrightarrow U_H \Leftrightarrow O_B \Leftrightarrow M \Leftrightarrow \left(U_H \triangleright \left(\left(H \Leftrightarrow \tilde{U}_H \Leftrightarrow O_B \Leftrightarrow M \right) \triangleleft \text{ADL} \right) \right) \right) \\ \triangleleft \left\langle \bigwedge_{i=1}^s f_{\text{ADL}}(S_i) \triangleleft \bigwedge_{i=1}^{n-1} fn_{\text{ADL}}(N_i) \triangleleft \left(N_n \Leftrightarrow U_{N_n} \Leftrightarrow \left(\left(\tilde{N}_n \Leftrightarrow \tilde{U}_{N_n} \right) \triangleleft \text{ADL} \right) \right) \right.$$

$$f_{\text{ADL}}(X) = X \Leftrightarrow U_X \Leftrightarrow \left(U_X \triangleright \left(\left(X \Leftrightarrow \tilde{U}_X \right) \triangleleft \text{ADL} \right) \right),$$

$S_1 \cdots S_s$ = surrounding non-namespace scopes,

$$fn_{\text{ADL}}(N) = N \Leftrightarrow U_N \Leftrightarrow \left(\left(N \Leftrightarrow U_N \right) \triangleright \left(\left(\tilde{N} \Leftrightarrow \tilde{U}_N \right) \triangleleft \text{ADL} \right) \right),$$

$N_1 \cdots N_n$ = surrounding namespaces,

H = scope where name binding is triggered from,

O_B = built-in operators,

M = member scope of operator's first argument,

U_X = using declarations in scope X ,

\tilde{X} = non-function (template) declarations in scope X , and

ADL = associated namespaces of reference's arguments.

Empty, since
operator is a
reserved keyword.



C++ Operators Scoping Rules with ADL

$$\left(H \Leftrightarrow U_H \Leftrightarrow O_B \Leftrightarrow M \Leftrightarrow \left(U_H \triangleright \left(\left(H \Leftrightarrow O_B \Leftrightarrow M \right) \triangleleft ADL \right) \right) \right) \\ \triangleleft \left(\bigwedge_{i=1}^s f_{ADL}(S_i) \triangleleft \left(\bigwedge_{i=1}^{n-1} fn_{ADL}(N_i) \triangleleft (N_n \Leftrightarrow U_{N_n} \Leftrightarrow ADL) \right) \right)$$

$$f_{ADL}(X) = X \Leftrightarrow U_X \Leftrightarrow (U_X \triangleright (X \triangleleft ADL)),$$

$S_1 \cdots S_s$ = surrounding non-namespace scopes,

$$fn_{ADL}(N) = N \Leftrightarrow U_N \Leftrightarrow \left((N \Leftrightarrow U_N) \triangleright ADL \right),$$

$N_1 \cdots N_n$ = surrounding namespaces,

H = scope where name binding is triggered from,

O_B = built-in operators,

M = member scope of operator's first argument,

U_X = using declarations in scope X , and

ADL = associated namespaces of reference's arguments.



C++ Operators Scoping Rules with ADL

$$\left(\mathbf{H} \Leftrightarrow \mathbf{U}_\mathbf{H} \Leftrightarrow \mathbf{O}_\mathbf{B} \Leftrightarrow \mathbf{M} \Leftrightarrow \left(\mathbf{U}_\mathbf{H} \triangleright \left(\left(\mathbf{H} \Leftrightarrow \mathbf{O}_\mathbf{B} \Leftrightarrow \mathbf{M} \right) \triangleleft \mathbf{ADL} \right) \right) \right) \\ \triangleleft \left\langle \bigwedge_{i=1}^s \mathbf{f}_{\mathbf{ADL}}(\mathbf{S}_i) \triangleleft \left(\left\langle \bigwedge_{i=1}^n (\mathbf{N}_i \Leftrightarrow \mathbf{U}_{\mathbf{N}_i}) \right\rangle \Leftrightarrow \mathbf{ADL} \right) \right.$$

$$\mathbf{f}_{\mathbf{ADL}}(\mathbf{X}) = \mathbf{X} \Leftrightarrow \mathbf{U}_\mathbf{X} \Leftrightarrow (\mathbf{U}_\mathbf{X} \triangleright (\mathbf{X} \triangleleft \mathbf{ADL})),$$

$\mathbf{S}_1 \cdots \mathbf{S}_s$ = surrounding non-namespace scopes,

$\mathbf{N}_1 \cdots \mathbf{N}_n$ = surrounding namespaces,

\mathbf{H} = scope where name binding is triggered from,

$\mathbf{O}_\mathbf{B}$ = built-in operators,

\mathbf{M} = member scope of operator's first argument,

$\mathbf{U}_\mathbf{X}$ = using declarations in scope \mathbf{X} , and

\mathbf{ADL} = associated namespaces of reference's arguments.



C++ Operators Scoping Rules

when scope **H** is an inner namespace scope

$$\left(\mathbf{H} \Leftrightarrow \mathbf{U}_\mathbf{H} \Leftrightarrow \mathbf{O}_\mathbf{B} \Leftrightarrow \mathbf{M} \Leftrightarrow \left(\left(\mathbf{H} \Leftrightarrow \mathbf{U}_\mathbf{H} \right) \triangleright \left(\left(\mathbf{O}_\mathbf{B} \Leftrightarrow \mathbf{M} \right) \triangleleft \mathbf{ADL} \right) \right) \right) \triangleleft \left(\bigvee_{i=1}^n (\mathbf{N}_i \Leftrightarrow \mathbf{U}_{\mathbf{N}_i}) \Leftrightarrow \mathbf{ADL} \right)$$

$\mathbf{N}_1 \cdots \mathbf{N}_n$ = surrounding namespaces,

\mathbf{H} = scope where name binding is triggered from,

$\mathbf{O}_\mathbf{B}$ = built-in operators,

\mathbf{M} = member scope of operator's first argument,

$\mathbf{U}_\mathbf{X}$ = using declarations in scope \mathbf{X} , and

\mathbf{ADL} = associated namespaces of reference's arguments.



C++ Operators Scoping Rules

when scope **H** is the outermost namespace scope

$$\left(\mathbf{H} \Leftrightarrow \mathbf{U}_{\mathbf{H}} \Leftrightarrow \mathbf{O}_{\mathbf{B}} \Leftrightarrow \mathbf{M} \Leftrightarrow \left(\left(\left(\mathbf{O}_{\mathbf{B}} \Leftrightarrow \mathbf{M} \right) \triangleleft \mathbf{ADL} \right) \right) \right)$$

H = scope where name binding is triggered from,

O_B = built-in operators,

M = member scope of operator's first argument,

U_X = using declarations in scope **X**, and

ADL = associated namespaces of reference's arguments.



2. Executing Scoping Rules

$$(\text{test} \Leftrightarrow \text{ns} \Leftrightarrow (\text{ns} \triangleright (\text{test} \triangleleft \text{adl}))) \triangleleft (:: \Leftrightarrow \text{adl})$$

1. Use **scope combinators**.

2. Instantiate the **Language** concept.

Bind

* **Best viable candidate, or**
* **error**

$$\text{bind} : \text{Ref} \times \text{Scope}_{\text{Ref}, \text{Decl}} \rightarrow (\text{Decl} + \text{Error})$$

```
void foo();  
namespace adl { struct X {};  
                void foo(X); }  
namespace ns { void foo(int); }  
void test(adl::X x) {  
    using ns::foo;  
    foo(x);  
}
```



Elementary Name Binding

Scopes as sets
of declarations.

```
lookup :: (...) => ref -> Set decl -> OverloadSet decl  
lookup ref decls = Set.filter (match ref) decls
```

```
resolve :: (...) => ref -> OverloadSet decl -> Maybe decl  
resolve ref decls = assess $ select_best_viable ref decls
```

```
assess :: (...) => BestViableSet decl -> Maybe decl  
assess decls = case (Set.elims decls) of  
    []      -> Nothing  
    [decl] -> (Just decl)  
    _       -> ambiguity decls
```



Elementary Name Binding

❖ What is the necessary minimal abstraction?

```
lookup :: (...) => ref -> Set decl -> OverloadSet decl  
lookup ref decls = Set.filter (match ref) decls
```

```
resolve :: (...) => ref -> OverloadSet decl -> Maybe decl  
resolve ref decls = assess $ select_best_viable ref decls
```

```
assess :: (...) => BestViableSet decl -> Maybe decl  
assess decls = case (Set.elims decls) of  
    []      -> Nothing  
    [decl] -> (Just decl)  
    _       -> ambiguity decls
```



An Abstraction of PLs

❖ The **Language** concept:

- `match`: How references match declarations.
- `select_best_viable`: How best viable candidates are determined.
- `ambiguity`: How to handle ambiguous results.



An Abstraction of PLs

❖ The **Ambiguity** concept:

- `ambiguity`: How to handle ambiguous results.

❖ The **Language** concept:

- `match`: How references match declarations.
- `select_best_viable`: How best viable candidates are determined.
- refines **Ambiguity**.



An Abstraction of PLs, Ext'd

❖ The **Ambiguity** concept:

- `ambiguity`: How to handle ambiguous results.

❖ The **Language** concept:

- `match`: How references match declarations.
- `select_best_viable`: How best viable candidates are determined.
- refines **Ambiguity**.

❖ The **Parameterized** concept:

- expresses changes in the bind environment during name binding.
- e.g., changes in **ambiguity** for a variant of *weak hiding*.



An Abstraction of PLs

```
class Ambiguity decl where  
  ambiguity :: ViableSet decl -> Maybe decl  
  ambiguity _ = Nothing -- default.
```

```
class Ambiguity decl => Language ref decl where  
  match :: ref -> decl -> bool  
  select_best_viable :: ref -> OverloadSet decl -> BestViableSet decl
```



An Abstraction of PLs, Ext'd

```
class Ambiguity decl where  
  ambiguity :: ViableSet decl -> Maybe decl  
  ambiguity _ = Nothing -- default.
```

```
class Ambiguity decl => Basic.Language ref decl where  
  match :: ref -> decl -> bool  
  select_best_viable :: ref -> OverloadSet decl -> BestViableSet decl
```

```
class (Parameterized decl, Basic.Language ref decl) =>  
  Language ref decl
```



Elementary Name Binding

```
lookup :: (Language ref decl) => ref -> Set decl -> OverloadSet decl  
lookup ref decls = Set.filter (match ref) decls
```

```
resolve :: (Language ref decl) => ref -> OverloadSet decl -> Maybe decl  
resolve ref decls = assess $ select_best_viable ref decls
```

```
assess :: (Ambiguity decl) => BestViableSet decl -> Maybe decl  
assess decls = case (Set.elims decls) of  
    []      -> Nothing  
    [decl] -> (Just decl)  
    _      -> ambiguity decls
```

❖ **assess** requires the **Ambiguity** concept only,
gets the meaning of ambiguity from **the language**.



Elementary Name Binding, Ext'd

```
resolve :: (Basic.Language ref decl) =>  
         BindEnv decl -> ref -> OverloadSet decl -> Maybe decl  
resolve env ref decls = assess env $ select_best_viable ref decls
```

```
assess :: BindEnv decl -> BestViableSet decl -> Maybe decl  
assess env decls = case (Set.elims decls) of  
    []      -> Nothing  
    [decl] -> (Just decl)  
    _       -> ambiguity env decls
```

❖ **assess** requires **no** concept,
gets the meaning of ambiguity from **the bind
environment**.



Applications

❖ Understanding current name binding mechanisms:

- Argument-dependent lookup.
- C++ operators.
- A cross-language analysis.

❖ Exploring concepts designs

- Understanding the current limitations.
- Exploring new solutions:
 - weak hiding,
 - 2-Stage name binding, and
 - parameterized weak hiding.

❖ Simplifying compiler designs.



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A cross language analysis

- ❖ Instantiate the **Language** concept with different languages, designs and kinds of references:
 - C++: function calls
 - **match** = have_**same_name** + **filter** out types.
 - **select_best_viable** = overload resolution (incl. viability check).
 - **ambiguity** = ambiguity_is_error.
 - C++: uses of types
 - **match** = have_**same_name**.
 - **select_best_viable** = viability check if single result, identity otherwise.
 - **ambiguity** = ambiguity_is_error.
 - C++: Multimethods proposal (cf. N2216)
 - **match** = (same as function calls).
 - **select_best_viable** = (same as function calls).
 - **ambiguity** = ambiguity_is_not_error.



C++ Multimethods (Std. Doc. N2216)

titled “Open Multi-Methods for C++”, Pirkelbauer et al., 2007

❖ Relaxes the rules for ambiguity,

- allowing ambiguous calls when all best viable candidates “have a **unique-base method** through which the call can be dispatched”.

```
struct X, Y, Z;
void foo(virtual X&, virtual Y&); // (1)
void foo(virtual Y&, virtual Y&); // (2) – ** unique-base method **
void foo(virtual Y&, virtual Z&); // (3)
struct XY : X, Y {}
struct YZ : Y, Z {}
void foo(virtual XY&, virtual Y&); // (4) – overrider for (1) and (2)
void foo(virtual Y&, virtual YZ&); // (5) – overrider for (2) and (3)

XY xy; YZ yz;
foo(xy,yz); // both (4) and (5) are equally viable matches,
            // w/ unique base (2)
```



A cross language analysis

- ❖ Instantiate the **Language** concept with different languages, designs and kinds of references:
 - Haskell: name uses, excluding type inference
 - `match =` `have_same_name.`
 - `select_best_viable =` `identity.`
 - `ambiguity =` `ambiguity_is_error.`
 - Haskell: name uses, including type inference
 - `match =` `have_same_name.`
 - `select_best_viable =` `viability check if single result, identity otherwise.`
 - `ambiguity =` `ambiguity_is_error.`
 - Alternatives have different implications with weak hiding or **Bind^{x2}**.



Applications

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- A cross-language analysis.

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Problem: Current Limitations

- ❖ Current scoping rules break seemingly valid codes.

The diagram shows a C++ code snippet with several annotations and callouts:

```
concept Foo2<typename P> =  
    requires (P a, P b) { foo(a, b); }  
  
void foo(int) { }  
  
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```

Annotations and Callouts:

- A blue box labeled `::` points to the namespace scope.
- A blue box labeled `Foo2<T>` points to the template parameter `Foo2 T`.
- A red callout box asks: "Should `Foo2<T>::foo()` shadow `::foo()`? How?"
- A red callout box asks: "Reject or accept the call `foo(1)`?"



Practical Examples

❖ STL: `rotate()` and `move()`

- Two function declarations, different number of parameters.
- Name lookup only finds one, type-check fails.

❖ Plenoptic photography: Image rendering

- Two function declarations, different parameter types.
- Name lookup only finds one, type-check fails.

❖ STL: `common_type`

- Two type (function) declarations, different type parameters.
- Name lookup only finds one, type-check fails.



STL: rotate() and move()

```
// Specialization adapted from latest release of libstdc++
template<RandomAccessIterator I>
    requires Permutable<I>          // has move(ValueType<I>&&)
I rotate (I first, I middle, I last) {
    ...
    if (__is_pod(ValueType<I>) && (middle - first == 1)) {
        ValueType<I> t = std::move(*p);          //
        std::move(p+1, p+last-first, p);          //
        *(p + last - first - 1) = std::move(t);    //
    }          // Ok w/ constrained std::move(), But...
    ...          // Does not use the constraints.
}
```



STL: rotate () and move ()

```
// A novice's approach
template<RandomAccessIterator I>
    requires Permutable<I>          // has move(ValueType<I>&&)
I rotate (I first, I middle, I last) {
    ...
    if (__is_pod(ValueType<I>) && (middle - first == 1)) {
        ValueType<I> t = move(*p);           // Ok
        move(p+1, p+last-first, p);          // ...
        *(p + last - first - 1) = move(t);    // Ok
    }    // Compile error... or... Possible ADL issues...
    ...
}
```



STL: rotate () and move ()

```
// The right way...
template<RandomAccessIterator I>
    requires Permutable<I> // has move(ValueType<I>&&)
I rotate (I first, I middle, I last) {
    ...
    if (__is_pod(ValueType<I>) && (middle - first == 1)) {
        ValueType<I> t = move(*p); // Ok
        std::move(p+1, p+last-first, p); // Ok
        *(p + last - first - 1) = move(t); // Ok
    } // Shadow or No Shadow... But...
    ... // Requires change in existing implementation.
}
```



STL: rotate () and move ()

```
// Update MoveWritable concept ...
template<RandomAccessIterator I>
    requires Permutable<I> // has std::move(ValueType<I>&&)
I rotate (I first, I middle, I last) {
    ...
    if (__is_pod(ValueType<I>) && (middle - first == 1)) {
        ValueType<I> t = std::move(*p);           // Ok
        std::move(p+1, p+last-first, p);          //
        *(p + last - first - 1) = std::move(t);    // Ok
    }        // Compile error... or... No ADL issues...
    ...      // Maybe OK to not shadow?
}           // std::move() is not a customization point.
```



Plenoptic photo.: Image rendering

```
template<typename PixelType>
struct Radiance {
    typedef typename
        boost::multi_array<PixelType, 4> RadianceType;
    RadianceType pixels;      // 4D array of pixels
    ...
    void Read(const string &ImageFile, ...);
    Radiance<PixelType> Render_Basic(...);
    Radiance<PixelType> Render_Blended(...);
    ... // Several rendering variants ...
    void Print(const string& OutputFile, ...);
    ...
}
```



Plenoptic photo.: Image rendering

```
template<typename PixelType>
Radiance<PixelType> Radiance<PixelType>::Render_Blended(...) {
    ...
    RadianceType Rendered(boost::extents[Iy][Ix][1][1]); ...
    for (...) { ... // for each image pixel
        PixelType pixel_avg; ...
        for (...) { ... // for each direction
            pixel_avg += pixels[ri][rj][rl][rm]; ... // integrate pixel
        } ...
        Rendered[i][j][0][0] = move(pixel_avg); ...
    }
    return move(new Radiance<PixelType>(move(Rendered), Ix, Iy));
}
```



Plenoptic photo.: Image rendering

```
template<typename PixelType>
using MultiArrayIterator =
    boost::multi_array<PixelType, 4>::iterator

// Add safety
template<Regular PixelType>
    requires IndirectlyCopyable<MultiArrayIterator<PixelType>,
                                   MultiArrayIterator<PixelType>>

struct Radiance {
    typedef typename
        boost::multi_array<PixelType, 4> RadianceType;
    ...
}
```



Plenoptic photo.: Image rendering

```
template<Regular PixelType> ...
Radiance<PixelType>&& Radiance<PixelType>::Render_Blended(...) {
    ...
    RadianceType Rendered(boost::extents[Iy][Ix][1][1]); ...
    for (...) { ... // for each image pixel
        PixelType pixel_avg; ...
        for (...) { ... // for each direction
            pixel_avg += pixels[ri][rj][rl][rm]; ... // integrate pixel
        } ...
        Rendered[i][j][0][0] = move(pixel_avg); ... // Ok.
    } // Compile error.
    return move(Radiance<PixelType>(move(Rendered), Ix, Iy));
}
```



STL: common_type

```
concept Common<typename T, typename U> =  
    requires { common_type<T, U>::type;      axiom (...) { ... } };  
  
template<typename T, typename U>  
requires Common<T, U> && ...  
void gen_func(...) {  
    common_type<T, U>::type ...           // Ok.  
    common_type<int>::type ...             // Error.  
    common_type<int, double>::type ...      // Error.  
    common_type<int, double, char>::type ... // Error.  
}
```



Solution: Weak Hiding

```
concept Foo2<typename P> =  
    requires (P a, P b) { foo(a, b); }
```

```
void foo(int) { }
```

Foo2<T>

```
template<Foo2 T>
```

```
void gen_func(T a, T b) {
```

```
foo(a, b);
```

```
foo(1);
```

}

```
Foo2<T>::foo()  
should weakly hide  
::foo()!
```

Accept the call foo(1)!

The Weak Hiding Scoping Rule

$$s_1 \Leftarrow s_2 = (resolve_{ref} \circ lookup_{ref}) s_1 ?$$

$$lookup_{ref} s_1 : lookup_{ref} s_2$$

```
concept Foo2<typename P> =  
    requires (P a, P b) { ... }
```

```
void foo(int) { }
```

Foo2<T>

```
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```

::

```
gen_func < P  
    < (Concept<T> <=> T) <::
```

Result:

Binds **foo()** to **::foo()**.



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Implementing Weak Hiding

$$s_1 \Leftarrow s_2 = (resolve_{ref} \circ lookup_{ref}) s_1 ?$$
$$lookup_{ref} s_1 : lookup_{ref} s_2$$

$$bind (ref, scope) = ((resolve\ ref) \circ (lookup\ ref))\ scope$$

❖ Implementation = **Two-Stage Name Binding (Bind^{x2})**

1. Bind with inner scope: **s1.**
2. Bind with outer scope: **s2.**

❖ Bind^{x2} repeats name binding under different contexts.



Bind^{x2} for C++ Concepts

1. Within restricted scope:

- up to the outermost restricted scope.
- Disables ADL and some qualified name lookups.

2. In surrounding scope:

- normal lookup – including ADL.

```
concept Foo2<typename P> =  
    requires (P a, P b) { ... }  
  
void foo(int) { }  
  
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```



Applications

❖ Understanding current name binding mechanisms:

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Inner Scope Ambiguity

when ambiguity IS an error

Current C++ Concepts:

Reject!

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b);  
}
```



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Inner Scope Ambiguity

when ambiguity IS an error

Proposed Extension:

Accept!

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b);  
}
```



Inner Scope Ambiguity

when ambiguity IS an error

False

Repeat bind

$$s_1 \Leftarrow s_2 = (resolve_{ref} \circ lookup_{ref}) s_1 ?$$
$$lookup_{ref} s_1 : lookup_{ref} s_2$$

$$bind (ref, scope) = ((resolve \ ref) \circ (lookup \ ref)) \ scope$$



Inner Scope Ambiguity

when ambiguity IS an error

False

Repeat bind

$$S_1 \Leftarrow S_2 = (resolve_{ref} \circ lookup_{ref}) S_1 ?$$
$$lookup_{ref} S_1 : lookup_{ref} S_2$$

Not Current C++ Concepts:
Reject or bind to unintended!

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b);  
}
```



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Inner Scope Ambiguity

when ambiguity IS an error

True

temporary change in ambiguity

Do not repeat bind

$S_1 \Leftarrow S_2 = (resolve_{ref} \circ (update \circ lookup_{ref}))$

$lookup_{ref} S_1 : lookup_{ref} S_2$

Current C++ Concepts:

Reject!

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b);  
}
```



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Inner Scope Ambiguity

when ambiguity IS NOT an error

True

Do not repeat bind

$$s_1 \Leftarrow s_2 = (resolve_{ref} \circ lookup_{ref}) s_1 ?$$

$lookup_{ref} s_1 : lookup_{ref} s_2$

Not Current C++ Concepts:
Accept!

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b);  
}
```



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Inner Scope Ambiguity

when ambiguity IS NOT an error

False

Repeat bind

$$S_1 \Leftarrow S_2 = (resolve_{ref} \circ lookup_{ref}) S_1 ?$$
$$lookup_{ref} S_1 : lookup_{ref} S_2$$

Not Current C++ Concepts:
Accept unintended ambiguity!

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b);  
}
```



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Ambiguity for C++ Concepts:

What is the most desirable?

1. Ambiguity IS an error, always?

**Rejects desirable, or
binds to undesirable**

2. Ambiguity IS NOT an error, always?

Accepts undesirable

3. A middle ground option?

**Proposed Extension:
Accept only desirable! (?)**

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b); foo(b);  
}
```



Ambiguity for C++ Concepts

Our proposed extension

- ❖ Ambiguity IS NOT an error, when in restricted scope.
 - Similar to multimethods proposal N2216.
- ❖ Ambiguity IS an error, otherwise.

Proposed Extension:
Accept only desirable!

```
concept Foo2<typename P> {  
    void foo(P, int);  
    void foo(int, P);  
}  
  
template<Foo2 T>  
void gen_func(T a, int b) {  
    foo(b, b); foo(b);  
}
```



Revisiting Hinnant's Example

❖ Cf. C++ standard documents:

- C++ Library Reflector message c++std-lib-20050
- N2576: *"Type-Soundness and Optimization in the Concepts Proposal"*, Douglas Gregor, March 2008.

❖ Question: **How to type-check constrained templates?**

```
template <class T>
    requires std::DefaultConstructible<T> && std::Addable<T> &&
    std::Convertible<T::result_type,T> && std::CopyConstructible<T>
void test() {
    T s1;
    T s2;
    T s3 = s1 + s2 + T() + T() + T();
}

int main() { test<string>(); }
```



Hinnant's Example: Problem

❖ The desired (optimized) output

- is observed with unconstrained templates, but
- is not observed with constrained templates.

```
auto concept Addable<typename T, typename U=T> {  
    typename result_type;  
    result_type operator+(T const&, U const&);  
}  
  
template <class T>  
    requires ... std::Addable<T> ...  
void test() {  
    T s1;  
    T s2;  
    T s3 = s1 + s2 + T() + T() + T();  
}  
  
int main() { test<string>(); }
```

Unconstrained:

lv string + lv string
rv string += rv string
rv string += rv string
rv string += rv string

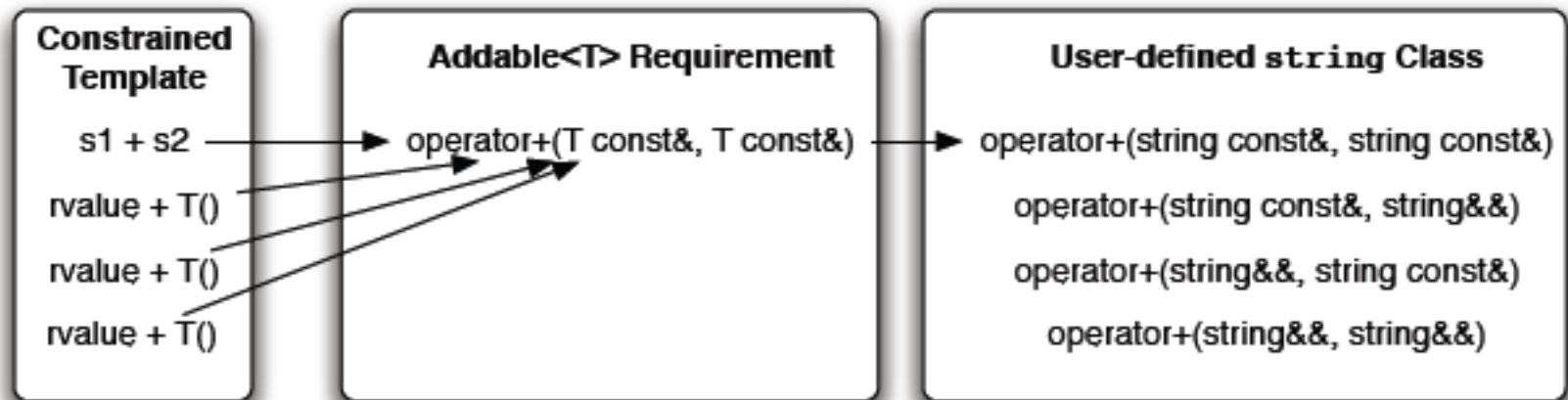
Constrained:

lv string + lv string
lv string + lv string
lv string + lv string
lv string + lv string



Hinnant's Example in Pictures

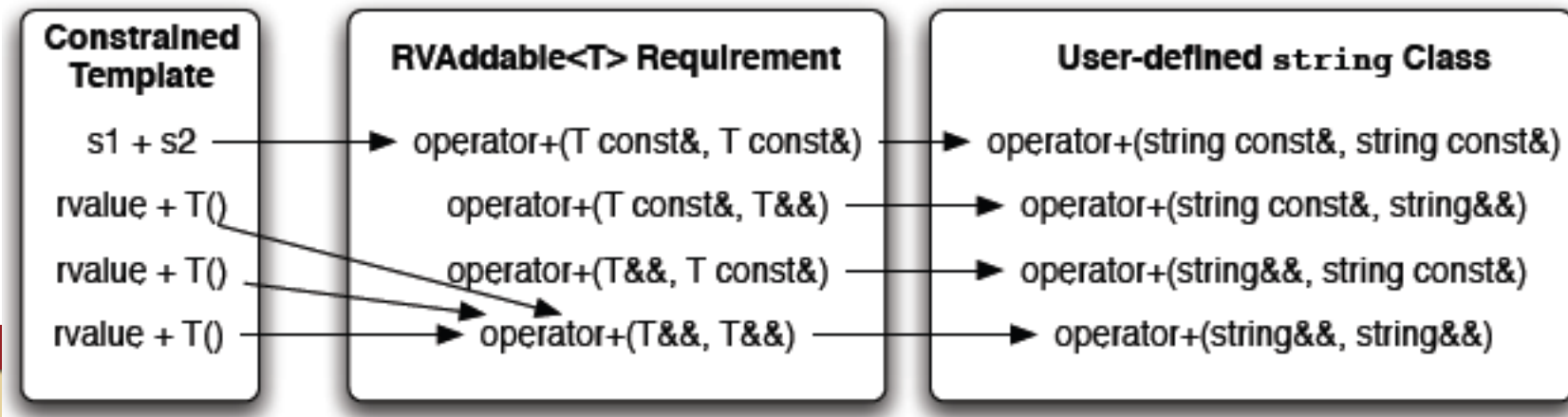
```
auto concept Addable<typename T, typename U = T> {  
    typename result_type;  
    result_type operator+(T const&, U const&);  
}
```



Hinnant's Example: Solution 1

Manual Introduction of Overloads

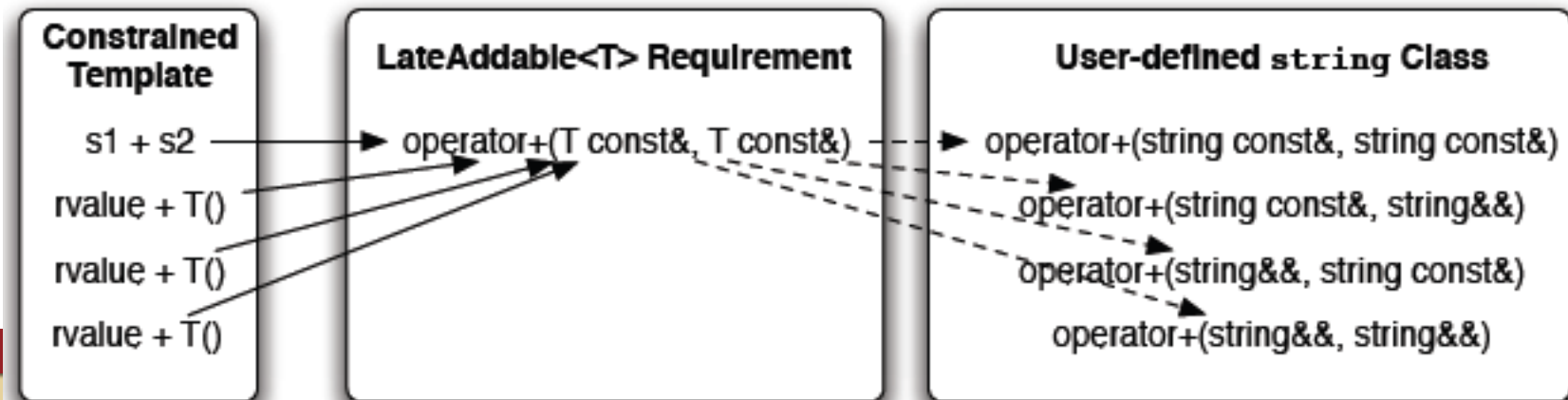
```
auto concept RVAddable<typename T, typename U = T> {  
    typename result_type;  
    result_type operator+(T const&, U const&);  
    result_type operator+(T const&, U&&);  
    result_type operator+(T&&, U const&);  
    result_type operator+(T&&, U&&);  
}
```



Hinnant's Example: Solution 2

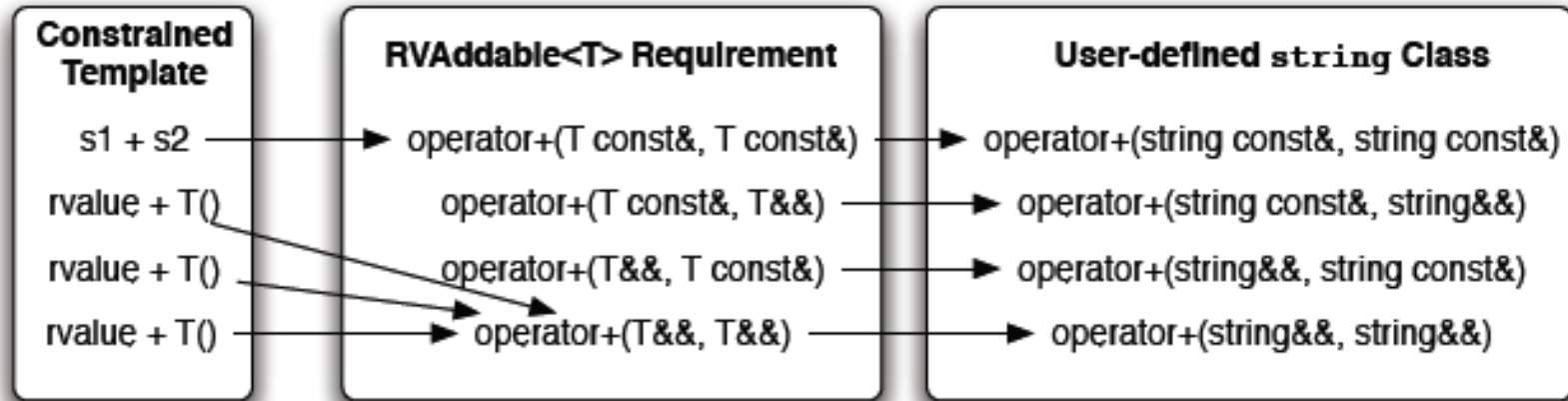
"Eliminating" Forwarding Functions

```
auto concept LateAddable<typename T, typename U = T> {  
    typename result_type;  
    result_type operator+(T const&, U const&);  
}  
concept_map LateAddable<string> {  
    typedef string result_type;  
    string operator+(T const& x, T const& y) { return x + y; }  
}
```

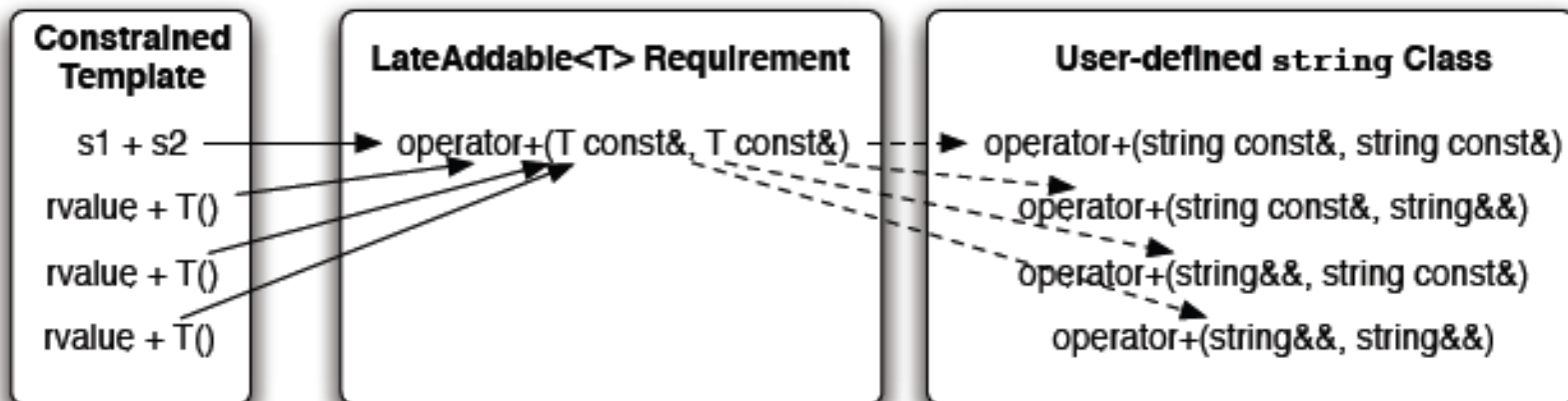


Revisiting the Proposed Solutions

1. Manual introduction of overloads:



2. “Eliminating” forwarding functions:



A “Sweet Middle” Alternative?

```
auto concept SweetAddable<typename T, typename U = T> {  
    typename result_type;  
    result_type operator+(T const&, U&&);  
    result_type operator+(T&&, U const&);  
}  
concept_map SweetAddable<string> {  
    typedef string result_type;  
    result_type operator+(T const& x, U&& y) { return x + y; }  
    result_type operator+(T&& x, U const& y) { return x + y; }  
}
```

Constrained Templates

s1 + T()

rvalue + T()

rvalue + s2

SweetAddable<T> Requirement

operator+(T const&, T&&)

operator+(T&&, T const&)

(ANY)

User-defined **string** class

operator+(string const&, string const&)

operator+(string const&, string&&)

operator+(string&&, string const&)

operator+(string&&, string&&)



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A “Sweet Middle” Example

Minimizing and clarifying template constraints

```
template <class T>
    requires ... std::SweetAddable<T> ...
void test() {
    T s1;
    T s2;
    T s3 = s1 + T() + T() + s2;
}

int main() { test<string>(); }
```

Constrained Templates

$s1 + T()$

$rvalue + T()$

$rvalue + s2$

SweetAddable<T> Requirement

$operator+(T\ const\&, T\&\&)$

$operator+(T\&\&, T\ const\&)$

(ANY)

User-defined `string` class

$operator+(string\ const\&, string\ const\&)$

$operator+(string\ const\&, string\&\&)$

$operator+(string\&\&, string\ const\&)$

$operator+(string\&\&, string\&\&)$



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Current C++ Concepts: Problem

- ❖ Clarity is penalized.
 - Minimal requirement specification leads to **ambiguity**,
 - which **is an error** in a constrained template definition,
 - even though **instantiation guarantees single matches**.
- ❖ Limited control over candidate sets.
 - Either **one or all** possible candidates.
- ❖ The “sweet middle” alternative is not allowed.



C++ Concepts: Proposed Extension

- ❖ **Ambiguity is not an error**, when in restricted scope.
- ❖ Ambiguity remains as defined in the language, otherwise.
- ➔ **Parameterize weak hiding** over the bind environment.



The Combinators, Revisited

- ❖ Weak hiding selectively updates lookup results

$$s_1 \Leftarrow s_2 = \left(\text{resolve}_{ref} \circ \left(\text{update} \circ \text{lookup}_{ref} \right) \right) s_1 ?$$
$$\left(\text{update} \circ \text{lookup}_{ref} \right) s_1 : \text{lookup}_{ref} s_2$$

- ❖ Define *update* based on a new concept: **Parameterized**.

```
update :: (Parameterized decl) => OverloadSet decl -> OverloadSet decl
update (os, env) = case get_env of
    Nothing          -> (os, env)
    (Just new_env)   -> (os, new_env)
```

Lookup results are coupled
with the bind environment.



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Parameterized example

```
class Parameterized decl where  
  get_env :: Maybe (ViableSet decl -> Maybe decl)  
  get_env = Nothing -- default, constant bind env.
```

Bind environment

```
class (Ambiguity decl, Parameterized decl,  
      Basic.Language ref decl) => Language ref decl
```

```
instance Ambiguity decl where  
  ambiguity = ambiguity_is_error  
  
instance Parameterized decl where  
  get_env = Just (ambiguity_is_not_error)
```

For C++ Concepts



Exploring Concepts Designs: Recap

- ❖ Introducing a new scoping rule: **Weak Hiding**.
- ❖ Introducing a new mechanism for name binding: **Bind^{x2}**.
- ❖ Main Motivation:
 - Software compatibility
 - In transition from unconstrained to constrained polymorphism.
 - e.g., a transition from C++ to ConceptC++.



Exploring Concepts Designs: Recap

- ❖ **Bind^{x2}** is an implementation of **weak hiding, optionally parameterized** (e.g., over the **bind environment**).
- ❖ Design variants: (Experimental)
 - Constant meaning of ambiguity: from **Ambiguity** concept.
 - Changing meaning of ambiguity: from **bind environment**.
 - e.g., Relaxing restrictions for name binding in restricted scope.



Applications

❖ Understanding current name binding mechanisms:

- Argument-dependent lookup.
- C++ Operators.
- A cross-language analysis.

❖ Exploring concepts designs

- Understanding the current limitations.
- Exploring new solutions:
 - weak hiding,
 - 2-Stage name binding, and
 - parameterized weak hiding.

❖ Simplifying compiler designs

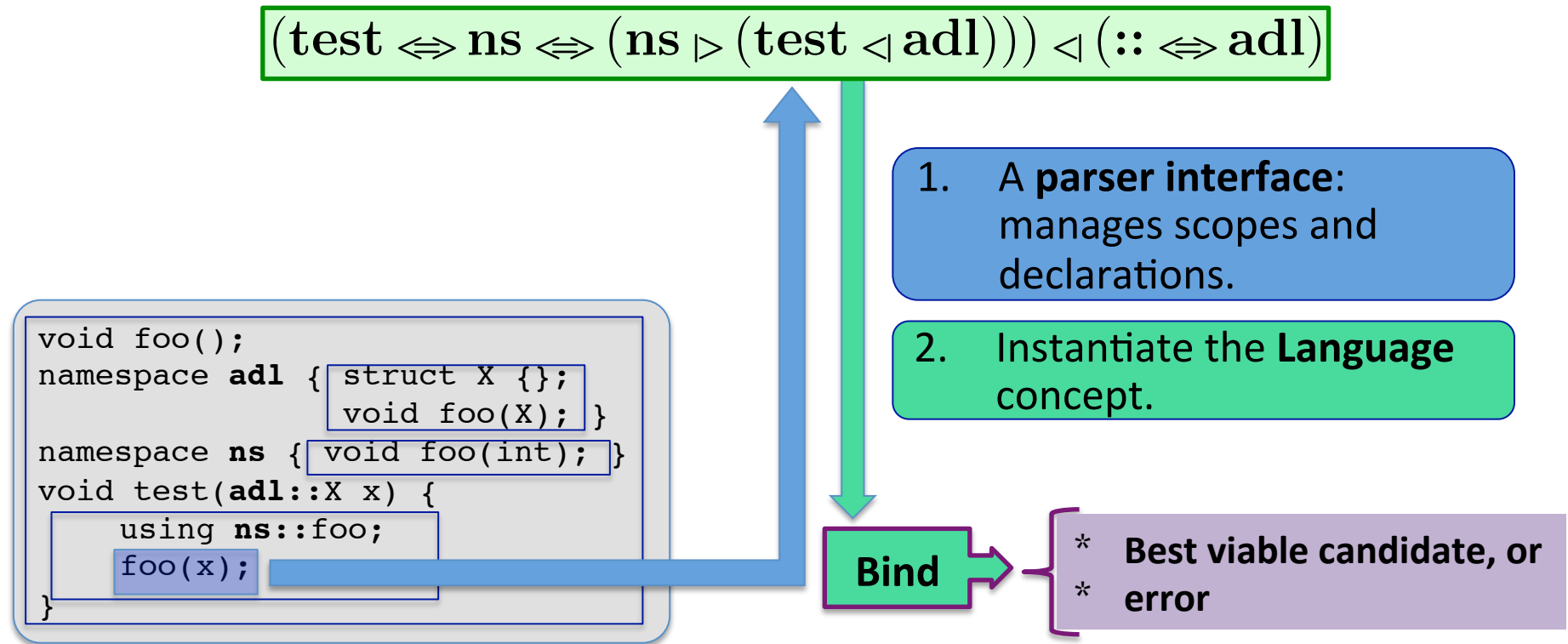


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Using our Framework in Compilers

Option #1: For all combinators, integrate directly into parsers



Using our Framework in Compilers

Option #2: for weak hiding only, a conservative integration

```
concept Foo2<typename P> =  
    requires (P a, P b) { ... }  
void foo(int) { }  
template<Foo2 T>  
void gen_func(T a, T b) {  
    foo(a, b);  
    foo(1);  
}
```

Bind^{x2}

* Best viable candidate, or
* error



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Prototypes Underway

[<https://github.iu.edu/lvoufo/BindIt>]

1. Parser interface for managing scopes:
 - Using Haskell's **Parsec** library and a **mini-C++** language.
 - Modifying existing compilers: **language-c**, **featherweight Java**, etc...
2. Executing **Bind^{x2}**:
 - Explored in **ConceptClang**.
3. Alternative compositional view of name binding
 - as composed of some **lookup_best** and **assess**.



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Recap

❖ Our name binding framework allows:

- expressing the scoping rules of a language, for a given reference, in terms of 3-4 **scope combinators**, and
- reasoning about the application of scoping rules generically, abstracting over the **Language** concept
 - incl. the **Ambiguity** and, optionally, **Parameterized** concepts.

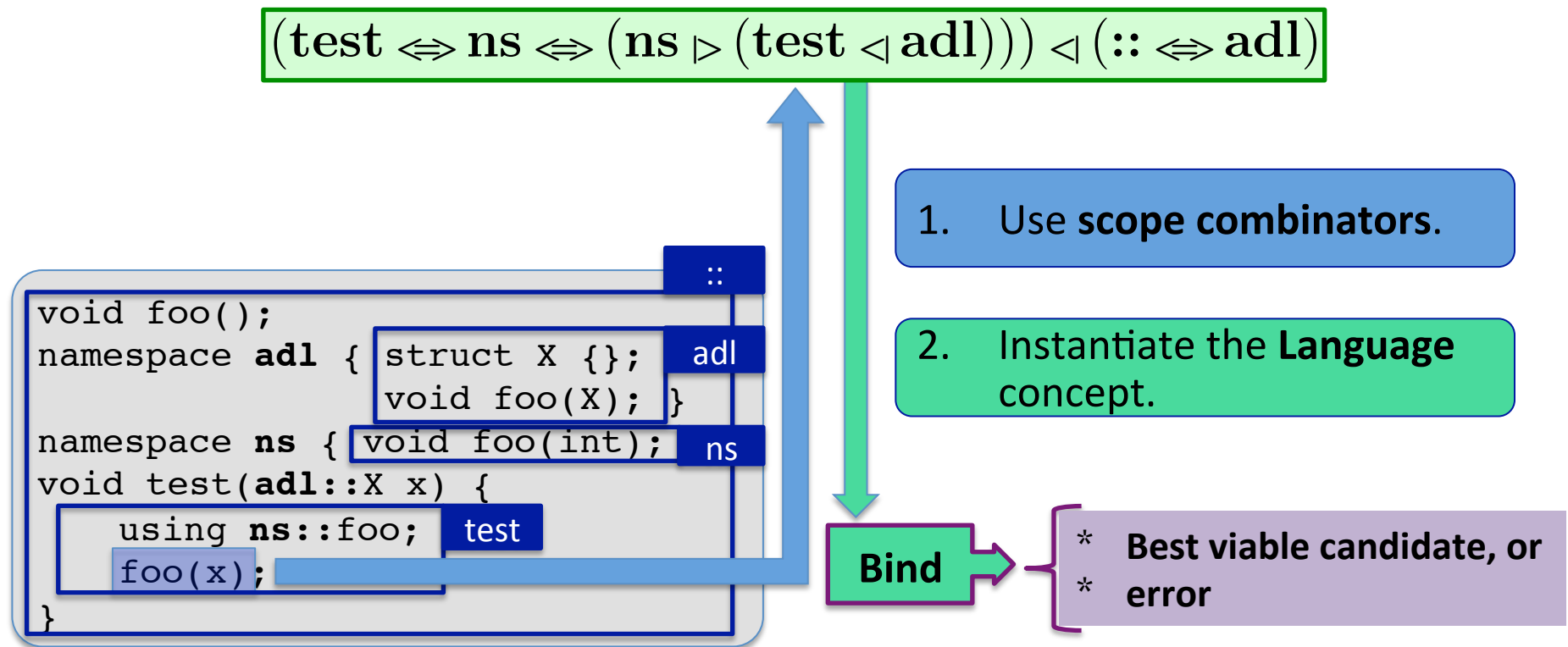
$\langle |$, $\langle \Rightarrow$, $| \rangle$, $\langle \Leftarrow$

❖ 2-stage name binding (**Bind^{x2}**):

- is an implementation of **weak hiding**, optionally parameterized by the bind environment, and
- preserves valid programs in transition from C++ to ConceptC++.



Thank You! Questions/Comments?



[<https://github.iu.edu/lvoufo/BindIt>]



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